

Lectures 20, 21: Bridges and switches

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314 S1T 2008

Repeaters to Routing

- How do we connect devices

- In the same room?

- In the next room?

- In the next building?

- In the next department

- In the next campus?

- In the next country?

- In the next continent?

Cabling limit

+ Admin boundary

+ Funding boundary

+ Political boundary

Distance, cost and numbers increase; must share costs and cables

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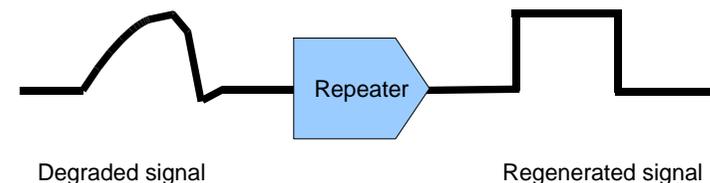
Topics (Shay 10.1 - 10.4)

- Repeaters
- Hubs
- Bridges
 - Spanning tree algorithm
- Ethernet switches
 - Virtual LANs
- The need for routing

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Repeaters

- A repeater simply reshapes, amplifies and sends on the signal
 - Useful when cable length is enough to degrade the signal, or to safely connect two cable segments



- An Ethernet repeater works in both directions

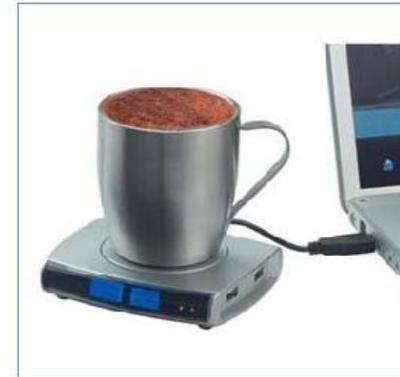
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Hubs

- A hub is (in the simplest case) simply a multiport repeater.
- A signal coming in on any port is reshaped, amplified, and sent out on all the other ports
- But *hub* is a very general word, so people sometimes speak of *bridging* or *switching* or even *routing* hubs.
- However, if it costs \$20 it's only a repeating hub.

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OK, maybe \$27.81...



This handy USB cup warmer plugs into your computer to keep your drink warm. Features include 4 port USB hub, when connected to computer the LCD blue back light is activated, clock, alarm and temperature.

from:
\$27.81 per unit

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Limits on hubs and repeaters

- Hubs and repeaters are essentially cable extenders.
 - They add slightly to transmission delay.
- In Ethernet, they form part of the collision domain calculation, as we saw earlier.

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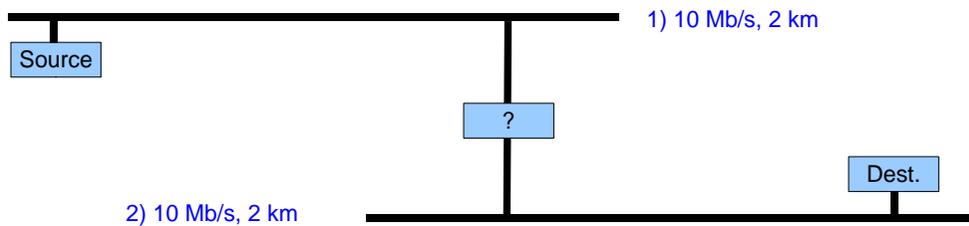
Advantages and disadvantages of hubs and repeaters

- + Nothing to configure.
- + Cheap and reliable.
- + Regenerate degraded signals.
- Don't avoid Ethernet contention/saturation problems on shared LANs.
- Don't split up the collision domain, therefore subject to distance limits.

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Bridges

- Although bridges are little used today, they are important to understand before we discuss switches.
- Consider a situation where we need to connect two Ethernets, but using a simple repeater would exceed the collision domain limit:



- What must we do to avoid joining the two collision domains into one?

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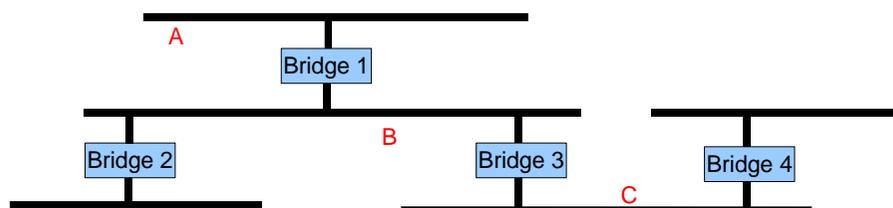
Bridges (2)

- The trick is to *read* the packet off one Ethernet and then *send* it on the other one quite independently.
 - The bridge "pretends" to be the destination on network 1 and then to be the source on network 2.
 - Reading a packet on Ethernet is a passive operation (the receiver sends absolutely nothing).
 - When sending the packet out again on network 2, the bridge uses the same source MAC address as the original sender on network 1, instead of its own MAC address.
 - Any station that sees the re-sent packet cannot tell that it came from a bridge.
 - Will any bits be different?
 - Does the bridge need to re-calculate the FCS?

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Bridges (3)

- How does the bridge know when to resend a packet?
 - In the simple case of two Ethernets, it could just resend everything it sees ("flooding").
 - But with several Ethernets and bridges, that gets too noisy, since every packet will be sent through every bridge to every other Ethernet.



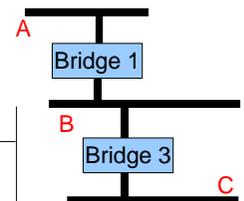
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Bridges (4)

- Conclusion: bridges must know where each the Ethernet station is, and only send packets in that direction.

- Each bridge needs a routing table. For Bridge 1:

From North		From South	
Dest. addr.	Send to:	Dest. addr.	Send to:
A	(Local)	A	North
B	South	B	(Local)
C	South	C	(Local)



- No need to send packets again locally

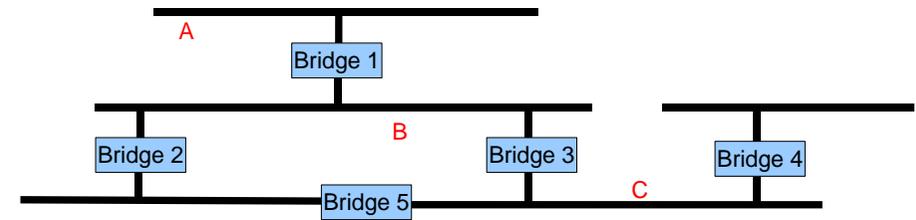
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How bridges learn

- At start-up, a bridge knows nothing (routing table is empty).
- Suppose it sees a packet on its North port, with destination B and source A.
- It now knows that A is to the north, so it can insert
A (Local) A North
- It still knows nothing of B, so the only choice is flooding: re-send the packet on its South port. All bridges do the same. (If B was in fact north, the re-sent packet would be wasted.)
- If it sees a packet with destination C and source A, it learns nothing new and continues to flood.
- When a reply comes from B to A on its South port, it inserts
B South B (Local)
- Try exercise 2 on Shay page 521

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Making the problem a little harder



- How do we prevent packets being sent twice or even in loops, when A sends to C?
 - Flooding and learning are not enough (both B2 and B3 will think that C is to the south...)
 - B2 needs to ignore packets for C on its North port.
 - The bridges will have to communicate with each other, to eliminate duplicate paths.

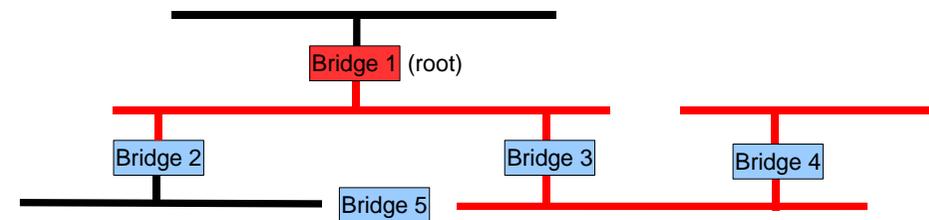
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The spanning tree algorithm in brief

- To avoid loops and duplicates we need a map of some kind. Bridges use a *spanning tree*.
 - (If you've done COMPSCI 220, you should know the theory behind this.)
- Bridges exchange management packets with each other.
- One bridge is "elected" as the *root*.
- Each bridge determines which port is closer to the root.
- The bridges on each LAN "elect" one bridge as the *designated bridge* which will send frames towards the root. Again, this is based on distance from the root.

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The spanning tree result



- Bridge 5 is idle, but there are no duplicates or loops.
 - Some shortest paths would use Bridge 5.
 - The spanning tree does *not* optimise paths.
- If Bridge 3 fails, the spanning tree will reconfigure and Bridge 5 will take over.
 - Routing tables and spanning tree must be timed out and refreshed periodically to deal with changes in the network

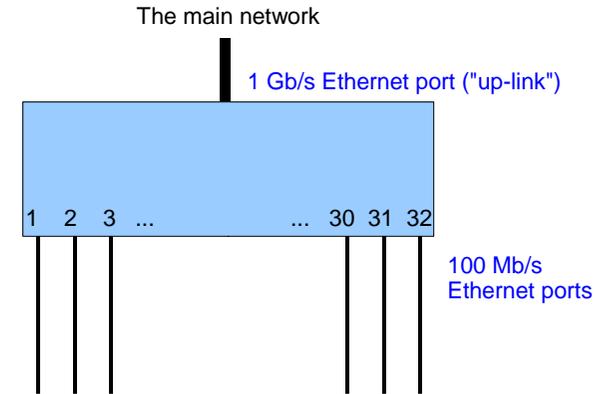
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Advantages and disadvantages of bridges

- + Self configuring. With spanning tree, any old network design will work - the easy way to split up a big LAN.
- + Fairly cheap and reliable.
- + Overcome distance limit, regenerate degraded frames.
- Don't avoid Ethernet contention/saturation problems on shared LANs.
- Suit old-fashioned shared cable better than modern UTP5 cabling.
- Notorious for allowing "broadcast storms" despite the spanning tree. (Bridges must forward multicast and broadcast packets.)

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Switches

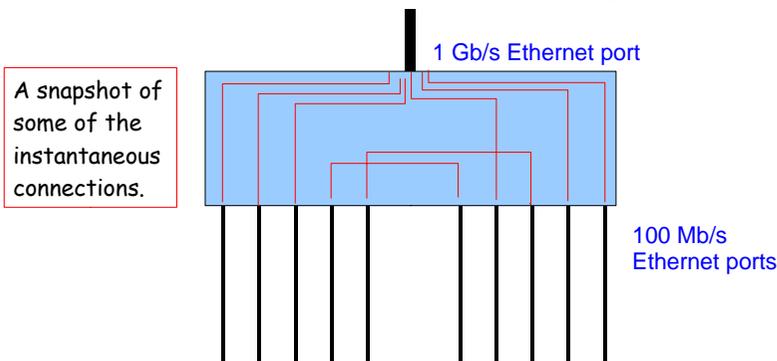


- Functionally, this switch will behave exactly like a 33 port bridge.
 - Routing table with 33 columns
 - Spanning tree

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Switch internals

- Switches vary widely in internal design (and price)
- The ideal switch can run full duplex at line speed on every port.
 - In our 33 port example, that would mean 32 100 Mb/s flows in and out. 10 each way could be using the 1 Gb/s port.



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Read the small print

- You can't tell from looking if a device is a switch or a hub.
- You can't tell from looking if a switch can run full speed on all ports, or if its internal design limits it to N frames per second or N simultaneous frames.
- In any case, 32x100 won't go into 1000. What happens if everyone sends to the "up-link" at the same time?
 - Does the switch queue excess packets, discard them, cause a collision, or apply some form of flow control?

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Example of small print

- 24 RJ-45, 10/100 Mbps Ethernet LAN ports
- Maximum Bandwidth:
 - Full duplex: 200 Mbps (for 100BASE-TX), 20 Mbps (for 10BASE-T)
 - Half duplex: 100 Mbps (for 100BASE-TX), 10 Mbps (for 10BASE-T)
- Operates at maximum forwarding rate and provides frame filtering and forwarding functions for each port
 - 14,880 packets per second for 10 Mbps ports
 - 148,800 packets per second for 100 Mbps ports
- Store-and-forward scheme to forward packets
- MAC address table: 4096 entries
- Packet buffer: 256 KB
- IEEE 802.3x Flow Control
 - Half-duplex: Back pressure
 - Full-duplex: Pause frame

USRobotics:
"24 Port 10/100 Mbps
Switch"

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Running the numbers

- "14,880 packets per second for 10 Mb/s ports"
- That means 10,000,000 / 14,880 bits per packet
- That's 672 bits per packet, or 84 bytes per packet.
- Ethernet packets are up to 1500 bytes, and more typically ~500 bytes.
- If we assume an average of 500 bytes/packet, that's 4000 bits/packet, which at 10 Mb/s means 10,000,000 / 4000 packets/s
- That's about 2500 packets per second, well below the switch capacity.
- "Packet buffer: 256 KB"
- That will hold about 500 typical packets. Suppose ten ports all fire packets at full speed towards the same port. 500 packets will arrive in 1/50 of a second (20 msec) but will take 200 msec to send on.
- Is the packet buffer big enough?
- In general, whether a switch is suitable depends on the expected traffic load. Obviously, switches with greater throughput cost more money.

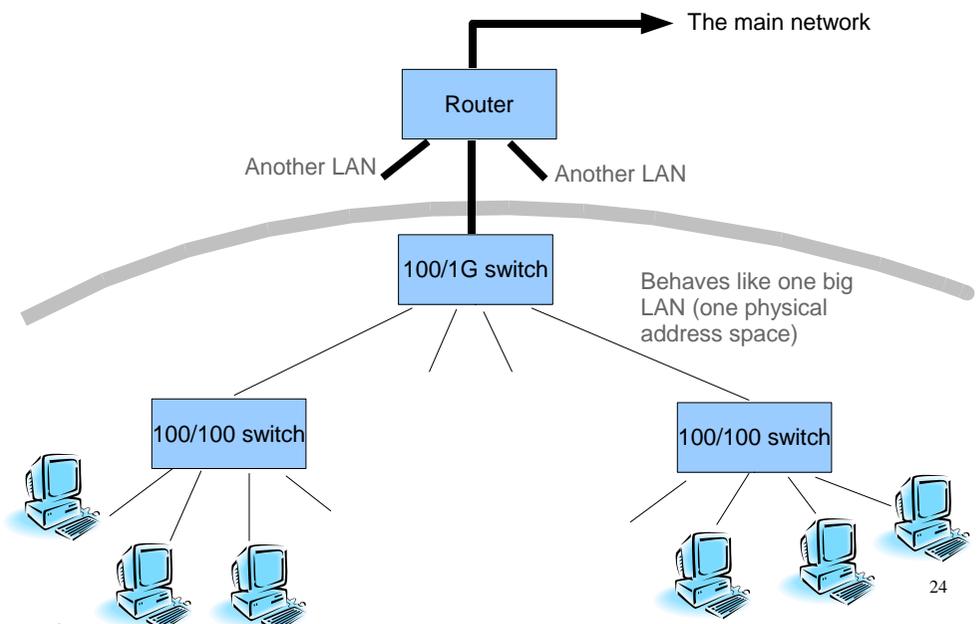
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Advantages and disadvantages of switches

- + Self configuring.
- + Reliable.
- + Overcome distance limit, regenerate degraded frames.
- + Avoid Ethernet contention/saturation problems on shared LANs or simple hubs.
- + Suit modern UTP5 cabling.
- Require systematic design and cable management.

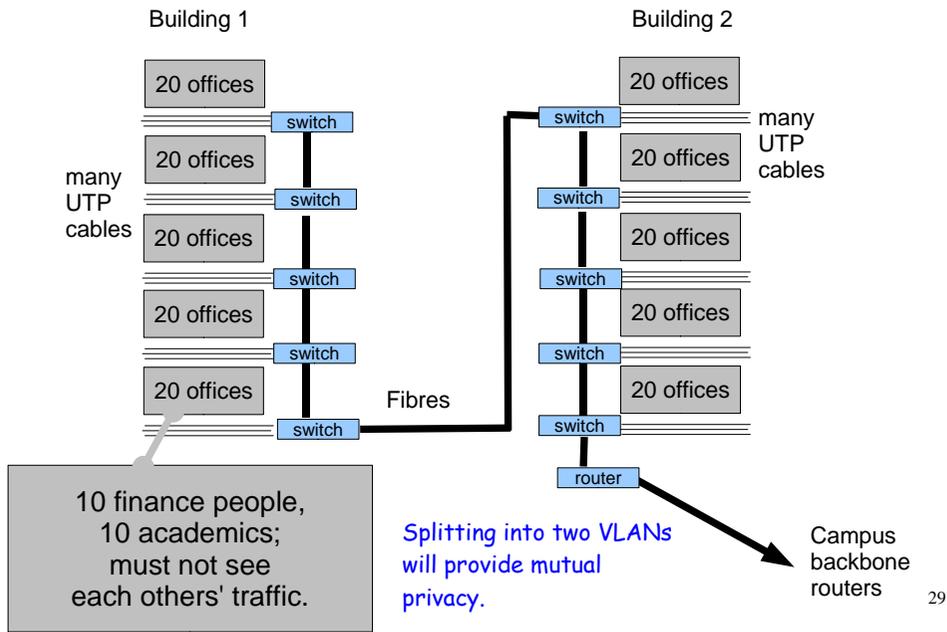
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Network for a large office or lab



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Two buildings, two populations



Repeaters to Routing

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- In the next continent?

Physical addressing:
Hubs, switches

Cabling limit

+ Admin boundary

+ Finding boundary

Logical addressing:
routers

+ Political boundary

Distance, cost and numbers increase; must share costs and cables, but separate administration